# Validated Numerics <br> a short introduction to rigorous computations 

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## Are floating point computations reliable?

Computing with the C/C++ single format

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Example 1: Repeated addition

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\begin{aligned}
& \sum_{i=1}^{10^{3}}\left\langle 10^{-3}\right\rangle=0.999990701675415 \\
& \sum_{i=1}^{10^{4}}\left\langle 10^{-4}\right\rangle=1.000053524971008
\end{aligned}
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Example 2: Order of summation

$$
\begin{aligned}
& 1+\frac{1}{2}+\frac{1}{3}+\cdots+\frac{1}{10^{6}}=14.357357 \\
& \frac{1}{10^{6}}+\cdots+\frac{1}{3}+\frac{1}{2}+1=14.392651
\end{aligned}
$$

## Are floating point computations reliable?

Given the point $(x, y)=(77617,33096)$, evaluate the function

$$
f(x, y)=333.75 y^{6}+x^{2}\left(11 x^{2} y^{2}-y^{6}-121 y^{4}-2\right)+5.5 y^{8}+x /(2 y)
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## IBM S/370 $(\beta=16)$ with FORTRAN:

| type | $p$ | $f(x, y)$ |
| :---: | :---: | :--- |
| REAL*4 | 24 | $1.172603 \ldots$ |
| REAL*8 | 53 | $1.1726039400531 \ldots$ |
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Pentium III $(\beta=2)$ with $C / C++(\mathrm{gcc} / \mathrm{g}++)$ :

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Correct answer: $-0.8273960599 \ldots$

## Round each partial result both ways

If $x, y \in \mathbb{F}$ and $\star \in\{+,-, \times, \div\}$, we can enclose the exact result in an interval:

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x \star y \in[\nabla(x \star y), \Delta(x \star y)] .
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## Question

How do we compute with intervals? And why, really?

## Arithmetic over $\mathbb{R} \mathbb{R}$

## Definition

If $\star$ is one of the operators,,$+- \times, \div$, and if $\boldsymbol{a}, \boldsymbol{b} \in \mathbb{R}$, then

$$
\boldsymbol{a} \star \boldsymbol{b}=\{a \star b: a \in \boldsymbol{a}, b \in \boldsymbol{b}\}
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Simple arithmetic

$$
\begin{aligned}
& \boldsymbol{a}+\boldsymbol{b}=[\underline{a}+\underline{b}, \overline{\boldsymbol{a}}+\overline{\boldsymbol{b}}] \\
& \boldsymbol{a}-\boldsymbol{b}=[\underline{\boldsymbol{a}}-\overline{\boldsymbol{b}}, \overline{\boldsymbol{a}}-\underline{\boldsymbol{b}}] \\
& \boldsymbol{a} \times \boldsymbol{b}=[\min \{\underline{\boldsymbol{a} b}, \underline{\boldsymbol{a}} \overline{\boldsymbol{b}}, \overline{\boldsymbol{a}} \underline{\underline{b}}, \overline{\boldsymbol{a}} \overline{\boldsymbol{b}}\}, \max \{\underline{\boldsymbol{a} b}, \underline{\boldsymbol{a}} \overline{\boldsymbol{b}}, \overline{\boldsymbol{a}} \underline{\boldsymbol{b}}, \overline{\boldsymbol{a}} \overline{\boldsymbol{b}}\}] \\
& \boldsymbol{a} \div \boldsymbol{b}=\boldsymbol{a} \times[1 / \overline{\boldsymbol{b}}, 1 / \underline{\boldsymbol{b}}], \quad \text { if } 0 \notin \boldsymbol{b} .
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On a computer we use directed rounding, e.g.

$$
\boldsymbol{a}+\boldsymbol{b}=[\nabla(\underline{\boldsymbol{a}} \oplus \underline{\boldsymbol{b}}), \triangle(\overline{\boldsymbol{a}} \oplus \overline{\boldsymbol{b}})] .
$$

Range enclosure
Extend a real-valued function $f$ to an interval-valued $F$ :

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R(f ; \boldsymbol{x})=\{f(x): x \in \boldsymbol{x}\} \subseteq F(\boldsymbol{x})
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$y \notin F(\boldsymbol{x})$ implies that $f(x) \neq y$ for all $x \in \boldsymbol{x}$.

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| $\sqrt{\boldsymbol{x}}$ | $=[\sqrt{\boldsymbol{x}}, \sqrt{\overline{\boldsymbol{x}}}]$ | if $0 \leq \underline{\boldsymbol{x}}$ |
| $\log \boldsymbol{x}$ | $=[\log \boldsymbol{x}, \log \overline{\boldsymbol{x}}]$ | if $0<\underline{\boldsymbol{x}}$ |
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Set $S^{+}=\{2 k \pi+\pi / 2: k \in \mathbb{Z}\}$ and $S^{-}=\{2 k \pi-\pi / 2: k \in \mathbb{Z}\}$. Then $\sin \boldsymbol{x}$ is given by

$$
\begin{cases}{[-1,1]} & : \text { if } \boldsymbol{x} \cap S^{-} \neq \emptyset \text { and } \boldsymbol{x} \cap S^{+} \neq \emptyset, \\ {[-1, \max \{\sin \underline{\boldsymbol{x}}, \sin \overline{\boldsymbol{x}}\}]} & : \text { if } \boldsymbol{x} \cap S^{-} \neq \emptyset \text { and } \boldsymbol{x} \cap S^{+}=\emptyset, \\ {[\min \{\sin \underline{\boldsymbol{x}}, \sin \overline{\boldsymbol{x}}\}, 1]} & : \text { if } \boldsymbol{x} \cap S^{-}=\emptyset \text { and } \boldsymbol{x} \cap S^{+} \neq \emptyset, \\ {[\min \{\sin \underline{\boldsymbol{x}}, \sin \overline{\boldsymbol{x}}\}, \max \{\sin \underline{\boldsymbol{x}}, \sin \overline{\boldsymbol{x}}\}]} & : \text { if } \boldsymbol{x} \cap S^{-}=\emptyset \text { and } \boldsymbol{x} \cap S^{+}=\emptyset .\end{cases}
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## Graph Enclosures

## A controlled discretization

We can now select and adapt the level of discretization

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Various levels of discretization of $f(x)=\cos ^{3} x+\sin x$.

## Solving non-linear equations



Solving non-linear equations


Consider everything. Keep what is good.
Avoid evil whenever you recognize it.
St. Paul, ca. 50 A.D. (The Bible, 1 Thess. 5:21-22)

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No solutions can be missed!

```
The code is transparent and natural
0 1 ~ f u n c t i o n ~ b i s e c t ( f c n N a m e , ~ X , ~ t o l ) ~
02 f = inline(fcnName);
03 if ( 0 <= f(X) ) % If f(X) contains zero...
04 if Diam(X) < tol % and the tolerance is met...
05 X % print the interval X.
0 6 ~ e l s e ~ \% ~ O t h e r w i s e , ~ d i v i d e ~ a n d ~ c o n q u e r .
07 bisect(fcnName, interval(Inf(X), Mid(X)), tol);
08 bisect(fcnName, interval(Mid(X), Sup(X)), tol);
09 end
10 end
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## Nice property

If $F$ is well-defined on the domain, the algorithm produces an enclosure of all zeros of $f$. [No existence is established, however.]

## Existence and uniqueness

Existence and uniqueness require fixed point theorems.

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## Brouwer's fixed point theorem

Let $B$ be homeomorhpic to the closed unit ball in $\mathbb{R}^{n}$. Then given any continuous mapping $f: B \rightarrow B$ there exists $x \in B$ such that $f(x)=x$.

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## Schauder's fixed point theorem

Let $X$ be a normed vector space, and let $K \subset X$ be a non-empty, compact, and convex set. Then given any continuous mapping $f: K \rightarrow K$ there exists $x \in K$ such that $f(x)=x$.

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## Banach's fixed point theorem

If $f$ is a contraction defined on a complete metric space $X$, then there exists a unique $x \in X$ such that $f(x)=x$.

## Theorem

Let $f \in C^{1}(\mathbb{R}, \mathbb{R})$, and set $\check{x}=\operatorname{mid}(\boldsymbol{x})$. We define

$$
N_{f}(\boldsymbol{x}) \stackrel{\text { def }}{=} N_{f}(\boldsymbol{x}, \check{x})=\check{x}-[D F(\boldsymbol{x})]^{-1} f(\check{x}) .
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Similar statements hold for the Krawczyk operator

$$
K_{f}(\boldsymbol{x}) \stackrel{\text { def }}{=} \check{x}-[D f(\check{x})]^{-1} f(\check{x})-\left(1-[D f(\check{x})]^{-1} F^{\prime}(\boldsymbol{x})\right)[-r, r],
$$

where we use the notation $r=\operatorname{rad}(\boldsymbol{x})$.

## Algorithm

Starting from an initial search region $\boldsymbol{x}_{0}$, we form the sequence

$$
\boldsymbol{x}_{i+1}=N_{f}\left(\boldsymbol{x}_{i}\right) \cap \boldsymbol{x}_{i} \quad i=0,1, \ldots .
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## Performance

If well-defined, this method is never worse than bisection, and it converges quadratically fast under mild conditions.

## Newton's method in $\mathbb{R}^{\mathbb{R}}$

## Example

Take $f(x)=-2.001+3 x-x^{3}$ and start with $\boldsymbol{x}_{0}=[-3,-3 / 2]$.

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X(0) = [-3.000000000000000,-1.5000000000000000]; rad = 7.50000e-01
X(1) = [-2.140015625000001,-1.546099999999996]; rad = 2.96958e-01
X(2) = [-2.140015625000001,-1.961277398284108]; rad = 8.93691e-02
X(3) = [-2.006849239640351,-1.995570580247208]; rad = 5.63933e-03
X(4) = [-2.000120104486270,-2.000103608530276]; rad = 8.24798e-06
X(5) = [-2.000111102890393,-2.000111102873815]; rad = 8.28893e-12
X(6) = [-2.000111102881727,-2.000111102881724]; rad = 1.55431e-15
X(7) = [-2.000111102881727,-2.000111102881724]; rad = 1.55431e-15
Finite convergence!
```

Unique root in -2.00011110288172 +- 1.555e-15

## Newton's method in $\mathbb{R}^{R}$

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## Stopping condition

Stop when no further improvement takes place.

## The Krawczyk method with bisection

When we have several zeros, we must bisect to isolate the zeros.

## Example

Take $f(x)=\sin \left(e^{x}+1\right)$ and start with $\boldsymbol{x}_{0}=[0,3]$.

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Domain : [0, 3]
Tolerance : 1e-10
Function calls : 71
    Unique zero in the interval 0.761549782880 [8890,9006]
    Unique zero in the interval 1.664529193[6825445,7060436]
    Unique zero in the interval 2.131177121086[2673,3558]
    Unique zero in the interval 2.4481018026567[773,801]
    Unique zero in the interval 2.68838906601606[36,68]
    Unique zero in the interval 2.8819786295709[728,1555]
```


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```


## Applications

Counting short periodic orbits for ODEs [Z. Galias] Measuring the stable regions of the quadratic map [D. Wilczak]

## Example

Draw the level-set defined by

$$
f(x, y)=\sin \left(\cos x^{2}+10 \sin y^{2}\right)-y \cos x=0
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restricted to the domain $[-5,5] \times[-5,5]$.

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MATLAB produces the following picture:


## Implicit curves (level sets)

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MATLAB produces the following picture:


According to the same m-file, the level set defined by $|f(x, y)|=0$, however, appears to be empty.

But this is the same set!!!

The (increasingly tight) set-valued enclosures in both cases are



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## Quadrature

Example (A bonus problem)
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A regular MATLAB session:
>> $q=q u a d(' \sin (x+\exp (x))$ ', 0,8$)$
$\mathrm{q}=$
0.251102722027180

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A regular MATLAB session:
>> $q=q u a d\left(' \sin (x+\exp (x))^{\prime}, 0,8\right)$
$\mathrm{q}=$
0.251102722027180

Using an adaptive validated integrator:
\$\$ ./adQuad 084 1e-4
Partitions: 8542
CPU time : 0.52 seconds
Integral : 0.347[3863144222905,4140198005782]

## Quadrature

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A regular MATLAB session:
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Using an adaptive validated integrator:
\$\$ ./adQuad 084 1e-4
Partitions: 8542
CPU time : 0.52 seconds
Integral : 0.347[3863144222905,4140198005782]
\$\$ ./adQuad 0820 1e-10
Partitions: 874
CPU time : 0.45 seconds
Integral : $0.3474001726[492276,652638]$

## Problem formulation

Given a finitely parametrized model function together with some (noisy) data, and a search region $\mathcal{P}$ in parameter space:

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- Instability: inverse problems can be extremely ill-conditioned.


## A statistical approach

Use a (weighted) least-squares approach to find the best parameter:

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\underset{p \in \mathcal{P}}{\operatorname{argmin}} \sum_{i=1}^{N} w_{i}\left|f\left(x_{i} ; p\right)-y_{i}\right|^{2} .
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## A set-valued approach

Locate nearby models that are consistent with nearby data:

$$
f(x ; p) \longrightarrow f(x ; \boldsymbol{p}) \quad\left(x_{i}, y_{i}\right) \longrightarrow\left(x_{i}, \boldsymbol{y}_{i}\right)
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Set-valued computations

## Points versus sets in parameter space

We move from the point-valued model function $f(x ; p)$ to the set-valued version $f(x ; \boldsymbol{p})$.

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Figure: (a) $p=0.15$, a point in $\mathcal{P}$. (b) $\boldsymbol{p}=[0.14,0.16]$, a subset of $\mathcal{P}$. The model function is $f(x ; p)=x e^{-p x}$, and 10 samples are shown.

## Strategy

Adaptively bisect the parameter space into sub-boxes: $\mathcal{P}=\cup_{j=1}^{K} \boldsymbol{p}_{j}$ and examine each $\boldsymbol{p}_{j}$ separately. [Good for parallelisation]

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(3) undetermined
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## Parameter estimation

## Example

Consider the model function

$$
f\left(x ; p_{1}, p_{2}\right)=5 e^{-p_{1} x}-4 \times 10^{-6} e^{-p_{2} x}
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with samples taken at $x=0,5 \ldots, 40$ using $p^{\sharp}=(0.11,-0.32)$.
Accepting a relative noise level of $90 \%$, we get the following set of consistent parameters:

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## Parameter estimation

Varying the relative noise levels between $10,20 \ldots, 90 \%$, we get the following indeterminate sets.


## Constraint propagation

Constraining the parameter/data space
We use set-valued constraint propagation to quickly discard inconsistent regions in the data and the parameter space.

This is done without bisection!

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f(2 ;[0,1])=2 e^{-2[0,1]}=2 e^{[-2,0]}=2\left[e^{-2}, 1\right]=\left[2 e^{-2}, 2\right] .
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This allows us to contract the data range according to

$$
\boldsymbol{y} \mapsto \boldsymbol{y} \cap f(x ; \boldsymbol{p})=[1,3] \cap\left[2 e^{-2}, 2\right]=[1,2] .
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Figure: The DAG representation of a forward sweep of $y=x e^{-p x}$, together with the corresponding code list.

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& n_{5}=n_{6} \div n_{1}=[1,2] \div 2=\left[\frac{1}{2}, 1\right] \\
& n_{4}=\log n_{5}=\log \left[\frac{1}{2}, 1\right]=[-\log 2,0] \\
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Note that, in one forward/backward sweep, we managed to exclude over $65 \%$ of the parameter domain, at the same time reducing the data uncertainty by $50 \%$.

## Mixed-effects models

We are given several data sets (trajectories) corresponding to $k$ different "individuals":

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\begin{array}{ll}
\text { individual }_{1}: & \left(x_{11}, y_{11}\right),\left(x_{12}, y_{12}\right), \ldots,\left(x_{1 N}, y_{1 N_{1}}\right) \\
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Some model parameters are equal (shared) for all individuals, and some are distinct.

- We need to consider all individuals simultaneously. Otherwise the number of unknown parameters may be too large.


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Search region:

$$
\mathcal{P}=([0,300],[0,9],[-1,0]) .
$$



Figure: Data inflation and contraction for the example. The graph of the model function for one subject (blue line). The data points are marked with red dots. The inflated data sets are shown as striped bars, and the re-contracted data as green bars.

Numerical results

|  | $N_{p}=1$ | $N_{p}=2$ |
| :--- | :---: | :---: |
| $\epsilon=0.01$ | $190.639(--)(0.010)$ | $193.141(19.6)(0.013)$ |
| $\epsilon=0.1$ | $194.139(--)(0.092)$ | $195.233(21.1)(0.097)$ |
| $\epsilon=0.2$ | $189.139(--)(0.190)$ | $193.437(20.3)(0.192)$ |
| $\epsilon=0.5$ | $167.226(--)(0.604)$ | $167.770(26.6)(0.589)$ |


|  | $N_{p}=5$ | $N_{p}=50$ |
| :--- | :---: | :---: |
| $\epsilon=0.01$ | $191.675(20.1)(0.014)$ | $191.239(20.1)(0.012)$ |
| $\epsilon=0.1$ | $192.954(21.4)(0.099)$ | $198.428(22.2)(0.110)$ |
| $\epsilon=0.2$ | $191.773(20.3)(0.203)$ | $197.580(23.6)(0.214)$ |
| $\epsilon=0.5$ | $164.656(23.9)(0.620)$ | $174.318(27.1)(0.618)$ |

Table: The results of four experiments for the example, each using 100 trial runs with $p_{1}=191.184$, and $\sigma=20.0$. For each pair $\left(\epsilon, N_{p}\right)$, we display the triple $\mu\left(p_{1}\right), \mu(\sigma)$, and $\mu(\epsilon)$ - the average estimates of the distribution parameters for $p_{1}$, and the data error.


Figure: The set of consistent parameters for two subjects from the example.

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- The relaxed problem is solved via set inversion.


## Further study...

## Interval Computations Web Page

http://www.cs.utep.edu/interval-comp

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INTLAB - INTerval LABoratory
http://www.ti3.tu-harburg.de/~rump/intlab/

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## CXSC - C eXtensions for Scientific Computation

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## CXSC - C eXtensions for Scientific Computation http://www.xsc.de/

CAPA - Computer-Aided Proofs in Analysis http://www.math.uu.se/~warwick/CAPA/

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## Validated Numerics:

A Short Introduction to Rigorous Computations

Warwick Tucker
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