- Functions and register conventions.
- Slide 1
- Stacks
- Implementing Stacks on the MIPS

•	As in high level languages , when programming in assembly
	language you should split up your program into smaller
	functions, that you can reuse.

Slide 2

• One of the key ideas with functions is that you can call them from any where and return back to where you called the function from.

• The MIPS processor has two instructions that enable you to call functions, jr and jal.

• Jump and link.

Slide 3

## jal label

Copies the address of the next instruction into the register **\$ra** (register 31) and then jumps to the address **label**.

• jr **\$register** jumps to the address in **\$register** most common use

jr \$ra

	str:	.data .asciiz "Hello mum!.\n" .text
	main:	.globl main #necessary for the assembler jal message jal message
Slide 4	message:	<pre>li \$v0,10 syscall #exit the program gracefully la \$a0,str li \$v0,4</pre>
		syscall #Magic to printhings on the screen. jr \$ra

• There are many way of passing values to functions, but there is a
convention that most programs on the MIPS follow.
• \$a0-\$a3 (registers 4 to 7) arguments 1-4 of a function.

• v0-v1 (registers 2 and 3) results of a function.

		li \$a0,10
		li \$a1,21
		li \$a3,31
		jal silly #Now the result of the function is is \$v0.
Slide 6	silly:	li \$v0,10
		syscall
		add \$t0,\$a0,\$a1
		sub \$v0,\$a3,\$t0
		jr \$ra

Slide 5

- On the previous slide in our function we needed an extra register to do part of a calculation.
- How do we know what registers to use?

As with function calls there is a convention.

- Slide 7
- \$s0-\$s7 the saved registers, these registers should be unchanged after a function call.
- **\$t0-\$t9** these are temporaries, are are not necessarily preserved across function calls.

So in the previous example it would of been a bad thing to use \$0 in the function silly.

• What happens if we run out of registers? What happens if we have to use **\$s0**?

Slide 8

• But where?

• We would have to save it.

Soon we will find a good place to store things.

		jal si	lly
		•	
		•	
		•	
<b>C11 1</b> 0	silly:	jal si	lly2
Slide 9		•	
		jr \$ra	

So we have to save **\$ra** as well.

- A stack is a data structure, at least two operations:
  - *push* put a value on the top of the stack
  - *pop* remove an item from the top of the stack.
- Slide 10
- The important thing about a stack is that it is a LIFO (Last in First Out) data structure. This is useful for nested functions.
  - You store your temporary data by pushing it onto the stack and restore things by popping things from it.

•	• The MIPS has no specialised <b>push</b> and <b>pop</b> instructions (Other processors do).
	• Instead the stack is implemented using the register \$sp (number 29), 1w and sw.

• Unless you are writing an operating system the register **\$sp** points to the top of the stack.

Slide 12

Slide 11

- $\bullet\,$  On the MIPS stacks grow downwards.
- You have to manipulate the value of the register **\$sp** and then use store and load.

To push the contents of register \$s0 onto the stack. Do the following:

```
addi $sp,$sp,-4
sw $s0,0($sp)
```

To pop the stop of the stack into register \$s0 do the following:

Slide 13 lw \$s0,0(\$sp) add \$sp,\$sp,4

Basic rules:

- Every thing you push onto the stack, you must pop from the stack.
- Never touch anything on the stack that does not belong to you.

	silly:	addi \$sp,\$sp,-4
		sw \$ra,0(\$sp)
<b>GIV1</b> 14		jal silly2
Slide 14		lw \$ra,0(\$sp)
		add \$sp,\$sp,4
		jr \$ra

How can we make the following code more efficient?

```
silly: addi $sp,$sp,-4
    sw $s0,0($sp)
    addi $sp,$sp,-4
    sw $ra,0($sp)
    jal silly2
    lw $ra,0($sp)
    addi $sp,$sp,4
    lw $s0,0($sp)
    addi $sp,$sp,4
    jr $ra
```

We have obeyed all the rules, but we are wasting some instructions. We don't need to add or subtract four twice, we could just add or subtract 8 and then change the loads and stores.

Slide 15

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General rule (applies to all programs you'll every write):

• Write the inefficient version once that is correct optimise.

result: .space 4 #the place for the result .text .globl main main: addi \$sp,\$sp,-4 #save the return address. lw \$a0, 5 Slide 17 jal fact la \$t0,result sw \$v0, 0(\$t0) lw \$ra,0(\$sp) addi \$sp,\$sp,4 jr \$ra

	fact:	addi \$sp,\$sp,-4
		sw \$ra,0(\$sp) #push \$ra on the stack
		#fact of 0 is 1
		bne \$a0,\$zero,not_zero
<b>G11 1</b> 4 6		#Set \$v0 to be 1
Slide 18		addi \$v0,\$zero,1
		#Restore \$ra from the stack
		lw \$ra,0(\$sp) #Read \$ra from the stack
		addi \$sp,\$sp,4 #restore the stack pointer.
		jr \$ra