

corbis.



Operating Systems - Spring 2009

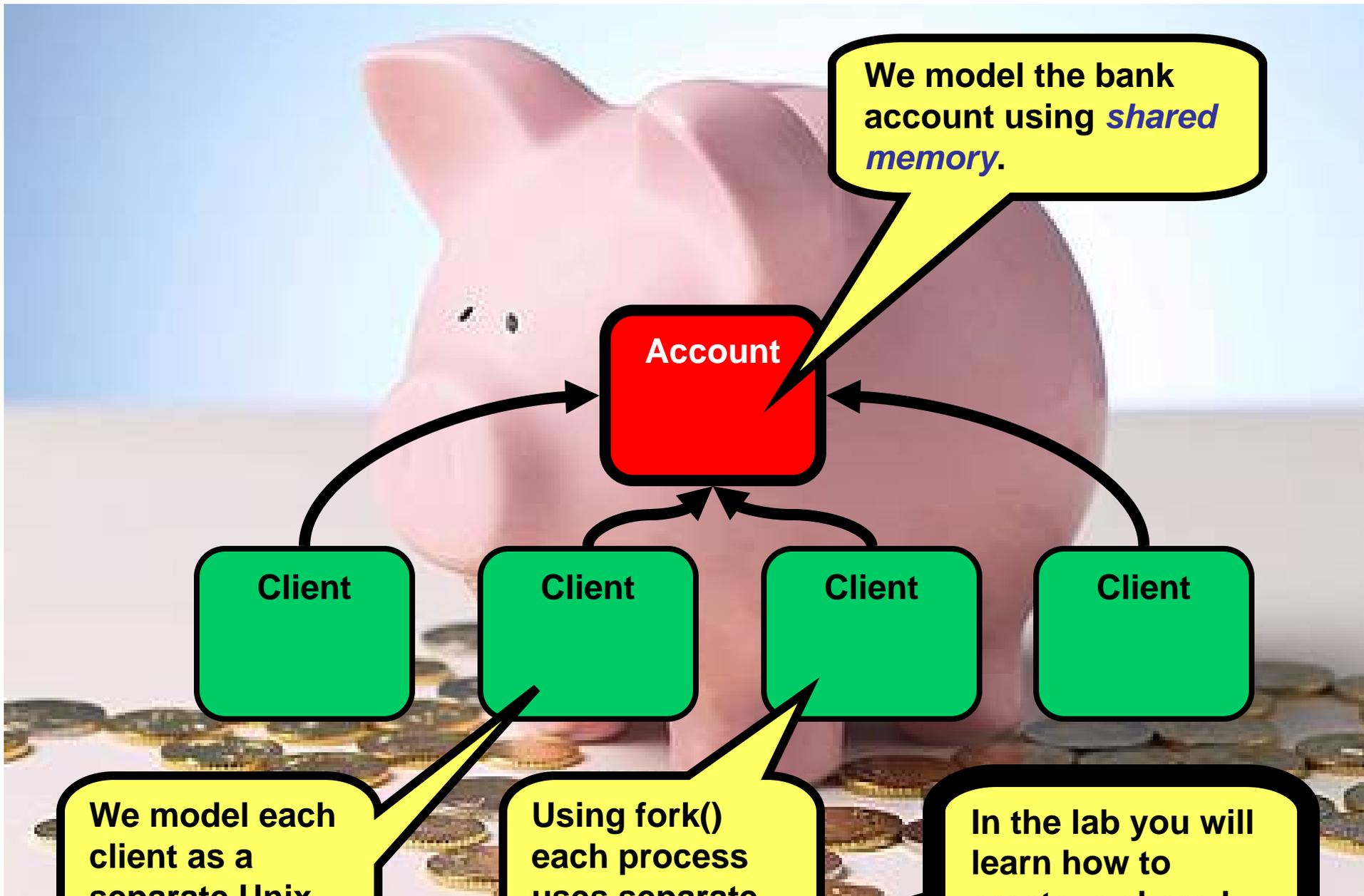
Assignment 2: Process Synchronization

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Assume we have one shared bank account.



All clients who has access to the account can make deposits.



We model the bank account using *shared memory*.

Account

Client

Client

Client

Client

We model each client as a separate Unix *process*.

Using `fork()` each process uses separate memory...

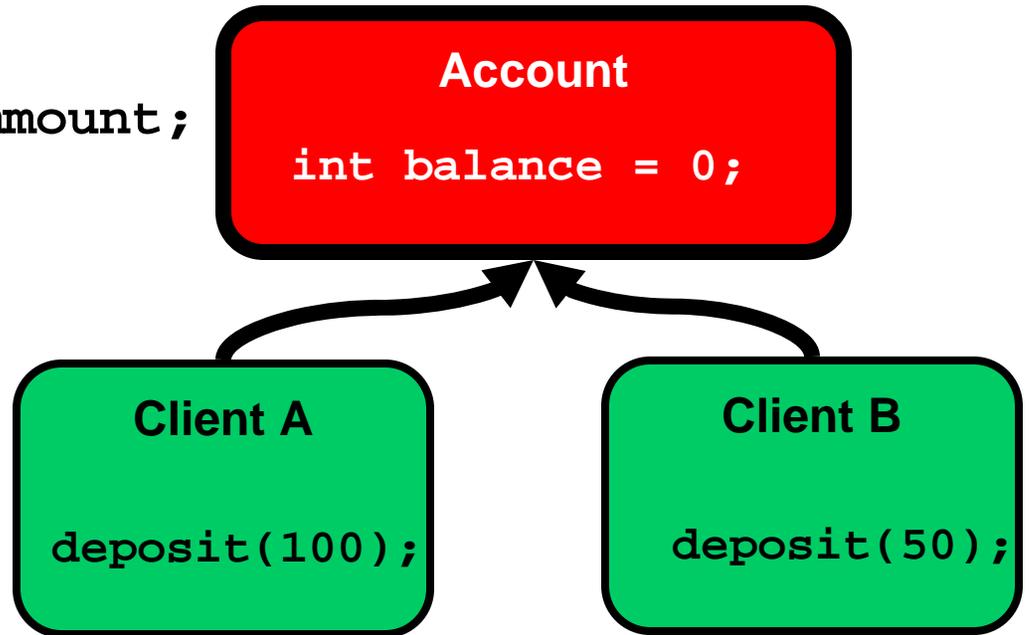
In the lab you will learn how to create a shared memory segment.

```

deposit(int ammount){
    balance = balance + ammount;
}

```

Remember: Incrementing the balance is not performed as one atomic operation by the hardware...



One possible scenario.

Client A	Client B	balance
Load (balance)		0
Add (balance, 100)		0
Store (balance)		100
	Load (balance)	100
	Add (balance, 50)	100
	Store (balance)	150

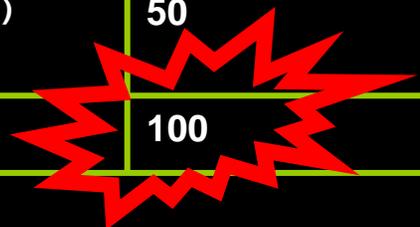
Another possible scenario.

Client A	Client B	balance
	Load (balance)	0
	Add (balance, 50)	0
	Store (balance)	50
Load (balance)		50
Add (balance, 100)		50
Store (balance)		150

Oops!

A *race condition* on the shared variable balance.

Client A	Client B	balance
	Load (balance)	0
	Add (balance, 50)	0
Load (balance)		0
Add (balance, 100)		0
	Store (balance)	50
Store (balance)		100



A blue, furry character with large white eyes and a wide black mouth is holding a large chocolate chip cookie in its right hand. The background is a blue gradient with light rays.

Must make sure only one process updates the account at the time...

Keep exactly one cookie in a jar.

Processes can look in the jar and *grap* the cookie if available.

If the cookie is absent – wait until another process *puts* it back.

Account
`int balance = 0;`

Client A
`Grap(cookie);`
`deposit(100);`
`Put(cookie);`

Client B
`Grap(cookie);`
`deposit(50);`
`Put(cookie);`



```
Account
int balance = 0;
```

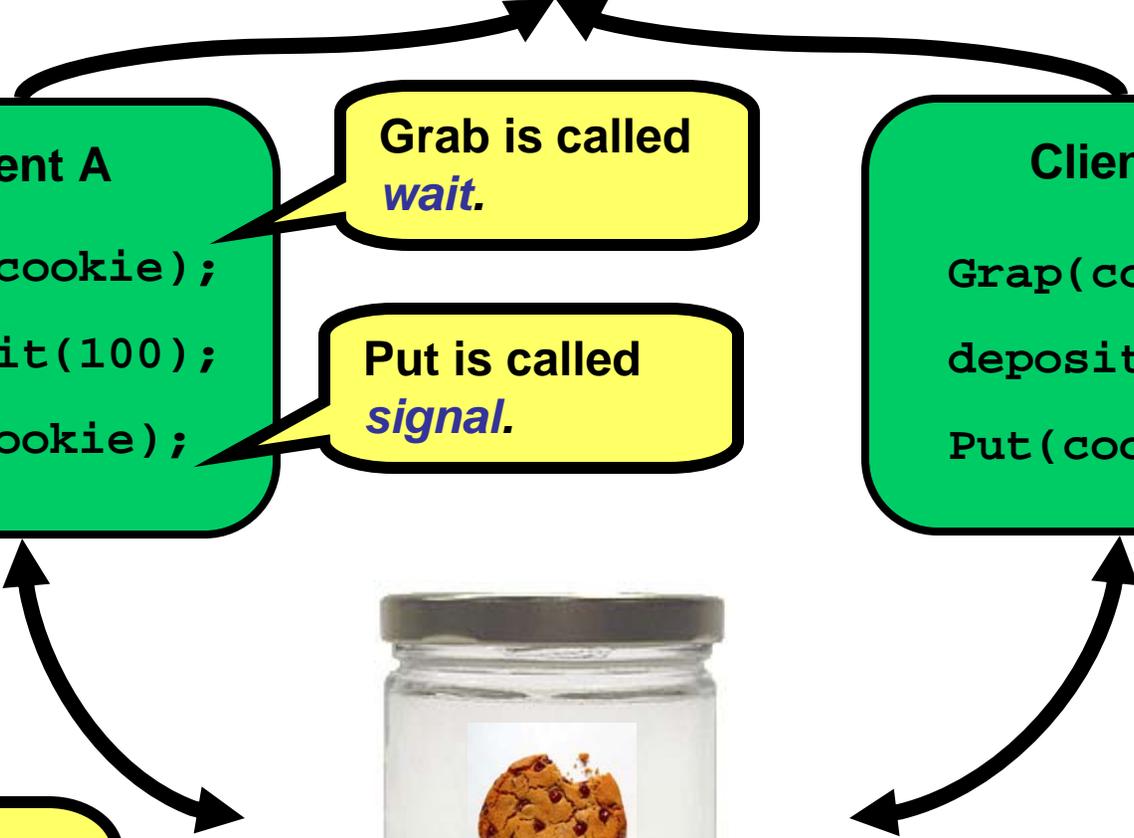
```
Client A
Grap(cookie);
deposit(100);
Put(cookie);
```

```
Client B
Grap(cookie);
deposit(50);
Put(cookie);
```

Grab is called *wait.*

Put is called *signal.*

The cookie jar is called a *binary semaphore.*





binary semaphore

counting semaphore



More than one "cookie" in the jar.

Signal: increments the semaphore counter.

Wait: if counter > 0 , decrement counter, otherwise wait.



In the lab you will learn how to use semaphores in C to synchronize processes.

Simplified API to
SystemV Semaphores.

```
#include "semaphore.h"

int n = 3;
Semaphore sem = Semaphore_create(n);

Semaphore_wait(sem);

Semaphore_signal(sem);

Semaphore_destroy(sem);
```

Create a new
counting
semaphore,
initialize
counter to n.

NOTE: you must destroy each semaphore
you create – they will not be deleted when
your program/process terminates.

**In the lab you will learn
how to use shared
memory to share data
between processes.**



Simplified API to SystemV
shared memory.

```
#include "shared_memory.h"

// Declare a handle to a shared memory segment.
Shared_memory shmId;

// Declare a pointer to data to be shared.
int *balance;

// Create the shared memory segment.
schmid = Shared_memory_create(sizeof(int));

// Must attach the segment to the
// process address space.
balance = (int*) Shared_memory_attach(shmId);

// Initialize data.
*balance = 1000;
```

NOTE: balance is a pointer to an int. Must use the *dereference operator* * to read/write the data stored at the pointer address.

```
// Don't forget to detach the segment from the  
// process address space when the process  
// is done using the segment.  
Shared_memory_detach(balance);
```

```
// Don't forget to destroy the shared segment  
// when all processes sharing the segment are  
// done using the segment.  
Shared_memory_destroy(shmid);
```

**NOTE: Shared memory segment
will not be deleted automatically
when a process exit.**