

PERFECT ROOT-OF-UNITY CODES WITH **PRIME-SIZE ALPHABET**

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Abstract

In this paper, Perfect Root-of-Unity Codes (PRUCs) with entries in $\alpha_p = \{x \in \mathbb{C} \mid x^p = 1\}$ where p is a prime are studied. A lower bound on the number of distinct phases in PRUCs over α_p is derived. We show that PRUCs of length $L \ge p(p-1)$ must use all phases in α_p . It is also shown that if there exists a PRUC of length L over α_p then p divides L. We derive equations (which we call principal equations) that give possible lengths of a PRUC over α_p together with their phase distribution.

Introduction

The Principal Equations

Now, for every $s \in \mathbb{Z}_p - \{0\}$, let us build the Φ_e matrix¹ as follows: by finding the locations of the entries $k_{l'}$ in the l^{th} column of Φ_x such that $k_{l'} \equiv k_l + s \pmod{p}$, we represent these locations in Φ_e by 1, and by 0 otherwise. Based on Corollary 2, for every $u \in \mathbb{Z}_L - \{0\}$, there exist exactly *m* distinct integers $\{l\}$ such that $k_{l+u} \equiv k_l + s \pmod{p}$. Therefore the Φ_e matrix has exactly m ones in each of its rows except for the first row which is all zero. This implies that Φ_e has m(mp-1) ones. On the other hand, the number of ones in Φ_e is equal to $\sum_{k=0}^{p-1} \mu_k \mu_{k+s}$ as it equals the number of all pairs with the property $k_{l+u} \equiv k_l + s \pmod{p}$. Therefore, all out-of-phase correlations $\{\sum_{k=0}^{p-1} \mu_k \mu_{k+s}\}_{s \in \mathbb{Z}_p - \{0\}}$ of the $\{\mu_k\}$ sequence are equal to m(mp-1).

Root-of-unity codes and their autocorrelation lags:

•
$$x = \{x_l\}_{l=0}^{L-1} = \left\{ e^{j\frac{2\pi}{p}k_l} \right\}_{l=0}^{L-1}$$

 L-length PRUC over α_p

• The periodic autocorrelation of x at lag $u \in \mathbb{Z}_L$ is defined as

$$R_{u} = \sum_{l=0}^{L-1} e^{j\frac{2\pi}{p}(k_{l}-k_{l+u})} = \begin{cases} L & u = 0\\ 0 & u \in \mathbb{Z}_{L} - \{0\} \end{cases}$$

Applications:

communication systems for Frequency-Hopping Spread-Spectrum Multiple-Access (FH/SSMA), Direct-Sequence Spread-Spectrum Multiple-Access (DS/SSMA), pulse compression for continuous-wave radars, fast startup equalization and channel estimation.

• This paper considers perfect unimodular codes in finite-alphabet case.

Phase Study

Theorem 1. If $\sum_{k=0}^{p-1} a_k e^{j\frac{2\pi}{p}k} = 0$ for some $a_k \in \mathbb{Z}$ and prime p, then all a_k must be identical.

• Let $r_k = \mu_k - m$, the principal equations:

$$\begin{cases} \sum_{k=0}^{p-1} r_k = 0\\ \sum_{k=0}^{p-1} r_k^2 = m(p-1)\\ \sum_{k=0}^{p-1} r_k r_{k+s} = -m, \qquad s \in \mathbb{Z}_p - \{0\} \end{cases}$$

• Geometrical study of the principal equations:

Let $r_0 = (r_0, \dots, r_{p-1})^T$ and also let r_k represent the circularly shifted version of r_0 by $k \in \mathbb{Z}_p$. The principal equations can be rephrased as follows over all vectors $\{r_k\}$:

> $\mathbf{1}_p^T \boldsymbol{r}_k = 0$ $\|\boldsymbol{r}_k\|_2 = \sqrt{m(p-1)}$ $\boldsymbol{r}_k^T \boldsymbol{r}_l = -m, \quad k \neq l$

The angle between each pair of vectors $\{(\mathbf{r}_k, \mathbf{r}_l)\}_{k \neq l}$ is

$$\theta = \arccos\left(\frac{\boldsymbol{r}_k^T \boldsymbol{r}_l}{\|\boldsymbol{r}_k\|_2 \|\boldsymbol{r}_l\|_2}\right) = \arccos\left(\frac{-1}{p-1}\right)$$



We note that the structure made by connecting all vertices pointed by $\{r_k\}$ is a known multi-dimensional object called a Regular Simplex.

Fig. 1. (a-c) Regular Simplexes in one, two and three dimensional space. In n dimensions they can be characterized with n + 1 vectors with the same l_2 -norm and also the same angle between them.

¹The dependency of $\{\Phi_e\}$ matrices on s is not explicitly shown for notational simplicity.

The Principal Equations: Solutions



• Lower bound on the number of distinct phases in the code:

Let Φ_x be the circulant matrix made from integer phases $\{k_l\}$ of the code

$$\Phi_x = \begin{pmatrix} k_0 & k_1 & \cdots & k_{mp-1} \\ k_{mp-1} & k_0 & \cdots & k_{mp-2} \\ \vdots & \vdots & \ddots & \vdots \\ k_1 & k_2 & \cdots & k_0 \end{pmatrix}$$

For the l^{th} column of Φ_x , consider the location of the entries which are equal to k_l $(l = 0, \dots, p-1)$. Considering these locations for all columns, we build an $mp \times mp$ equivalence matrix Φ_e whose entries in the mentioned locations are 1; otherwise they are 0.

- If μ_k represents the number of times that $e^{j\frac{2\pi}{p}k}$ occurs in the sequence then $\sum_{k=0}^{p-1} \mu_k = mp$.
- Based on Corollary 2, all rows of Φ_e have exactly m ones except the first row whose all entries are one. \square there are mp + m(mp - 1) ones in Φ_e .
- On the other hand, since every integer phase $k \in \mathbb{Z}_p$ gives μ_k columns with μ_k ones in each

Analytical solutions:

As an example for using the regular simplex to solve the principal equations we study the case of p = 3: for 3-phase perfect codes, the $\{r_k\}$ make a two dimensional regular simplex orthogonal to $\mathbf{1}_3$, which has 3 vectors and each two of them have an angle of $\frac{2\pi}{3}$. Let: R_{1_3} be the unitary rotation matrix which maps 1_3 to $\sqrt{3}e_3^{(3)}$ and $\mathbf{r}'_k = \sqrt{2m} \left(\cos \left(\frac{2k\pi}{3} + \psi \right), \sin \left(\frac{2k\pi}{3} + \psi \right), 0 \right)^T$ for $k \in \mathbb{Z}_3, \psi \in [0, 2\pi)$.



	$oldsymbol{r}_k \in \mathbb{Z}^3$ is of the form	
	$\left(\frac{\sqrt{2}}{2}\cos\left(\frac{2k\pi}{3}+\psi\right)-\frac{\sqrt{6}}{6}\sin\left(\frac{2k\pi}{3}+\psi\right)\right)$	
$\sqrt{2m}$	$\frac{\sqrt{6}}{3}\sin\left(\frac{2k\pi}{3}+\psi\right)$	
	$\left(-\frac{\sqrt{2}}{2}\cos\left(\frac{2k\pi}{3}+\psi\right)-\frac{\sqrt{6}}{6}\sin\left(\frac{2k\pi}{3}+\psi\right)\right)$	

For k = 0,

 $h_1 = 2\sqrt{\frac{m}{3}}\sin\psi$ (which is the second entry of r_0)

 $h_2 = 2\sqrt{m}\cos\psi$ (which is the difference between the first and the third entry of r_0

both must be integers.

Therefore,

the code length must be of the form

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Computational solutions:					
p	Length (L)	Phase distribution $\{\mu_k\}_{k=0}^{p-1}$			
	5	(2, 1, 2, 0, 0) (2, 2, 0, 1, 0)			
	20	(6, 4, 6, 2, 2) $(6, 6, 2, 4, 2)$			
	25	(6, 6, 6, 6, 1) $(9, 4, 4, 4, 4)$			
5	45	(12, 9, 12, 6, 6) $(12, 12, 6, 9, 6)$			
	55	(16, 12, 10, 10, 7) $(16, 10, 12, 7, 10)$			
		(15, 6, 10, 12, 12) $(15, 12, 6, 12, 10)$			
	80	(20, 16, 20, 12, 12) $(20, 20, 12, 16, 12)$			
	100	(28, 18, 18, 18, 18) $(22, 22, 22, 22, 12)$			
	7	(2, 2, 1, 0, 0, 2, 0)			
	28	(6, 6, 4, 2, 2, 6, 2) $(7, 5, 5, 2, 5, 2, 2)$			
7	49	(13, 6, 6, 6, 6, 6, 6) $(8, 8, 8, 8, 8, 8, 1)$			
	56	(11, 10, 10, 5, 10, 5, 5) $(11, 11, 6, 11, 6, 6, 5)$			
1	0.0				

of them, the number of ones in Φ_e is equal to $\sum_{k=0}^{p-1} \mu_k^2$ and therefore $\sum_{k=0}^{p-1} \mu_k^2 = mp + m(mp-1)$.

• Now let us assume that t of $\{\mu_k\}$ are nonzero. From the Cauchy-Schwarz inequality we have



Corollary 3. For $m \ge p - 1$, all phase values must be used in a PRUC.





Table 1. All possible lengths (less than or equal to 100) of PRUCs over α_p for p = 5 and 7 together with phase distributions.

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