

# On the Tractability of Digraph-Based Task Models

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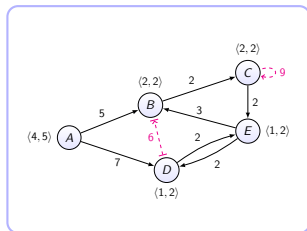
Joint work with Pontus Ekberg, Nan Guan and Wang Yi

# Analysis of Abstract Models

- Model hard real-time systems
  - ▶ Analysis: Guarantee deadlines
  - ▶ *Expressiveness* of models?
  - ▶ *Efficiency* of analysis?

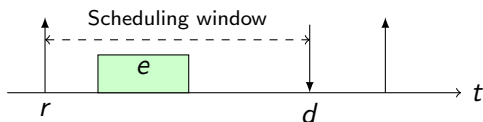
## Question

- How expressive can a model be?
- ... with a *tractable* feasibility test?



# Context: Real-Time Task Models

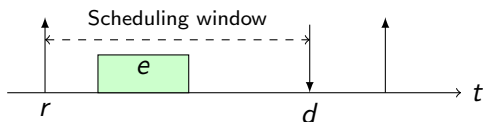
- System is composed of *tasks*, releasing *jobs*  $J = (r, e, d)$ 
  - ▶ Release time  $r$
  - ▶ Worst-case execution time  $e$
  - ▶ Deadline  $d$



- Feasibility: *Can we schedule s.t. all jobs meet their deadlines?*
  - In this work:
    - ▶ Preemptive schedules
    - ▶ On uniprocessors
    - ▶ Independent jobs
- } In this setting: EDF is optimal.  
*Feasible*  $\Leftrightarrow$  *Schedulable with EDF*

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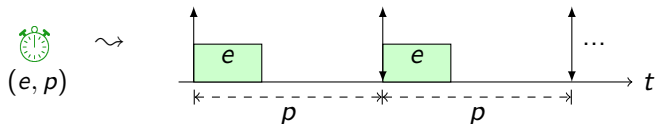


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# The Liu and Layland (L&L) Task Model

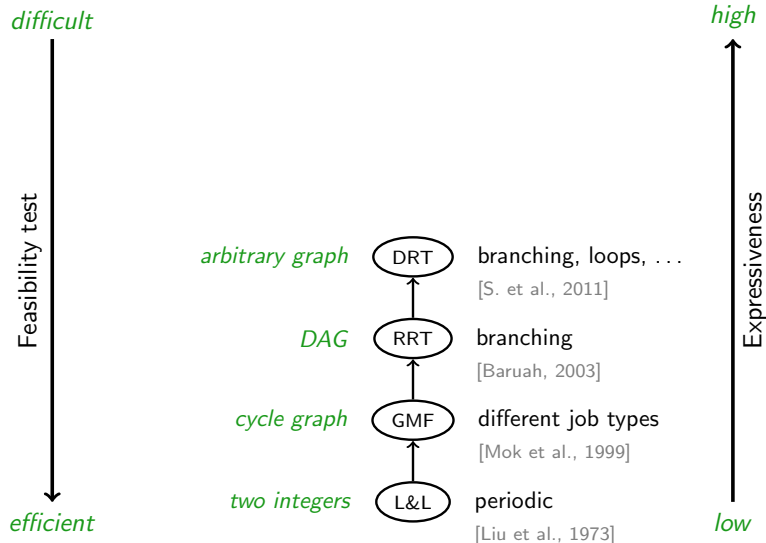
(Liu and Layland, 1973)

- Tasks are *periodic*
  - ▶ Job WCET  $e$
  - ▶ Minimum inter-release delay  $p$  (implicit deadline)

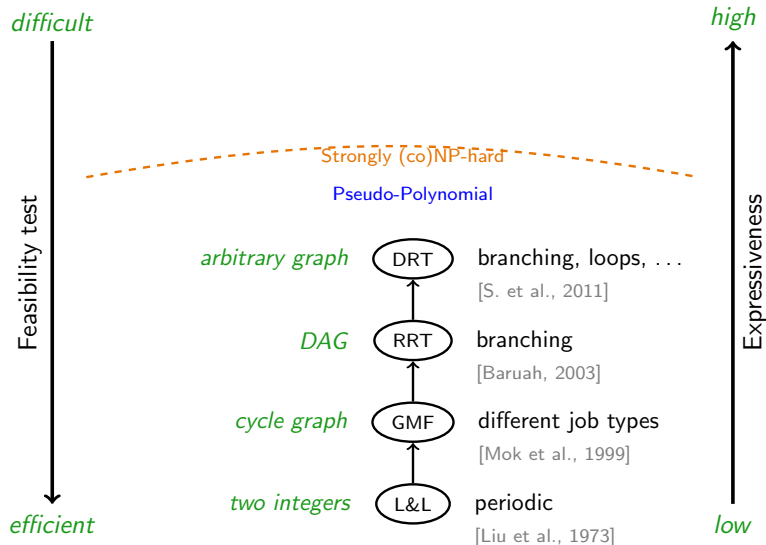


- Advantages: Well-known model; *efficient* schedulability test
- Disadvantage: Very *limited expressiveness*

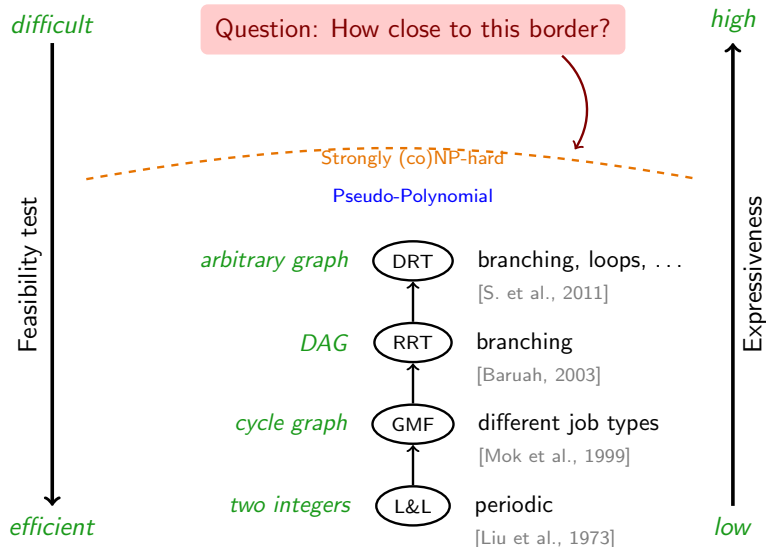
# Hierarchy of Models



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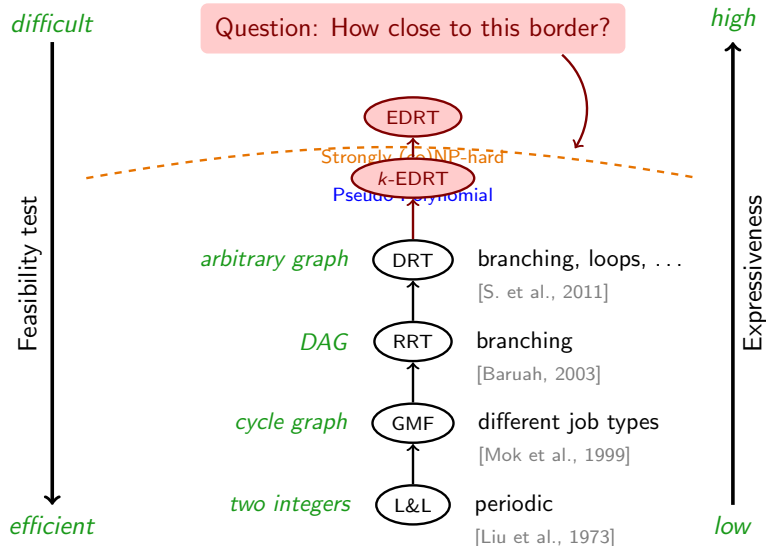


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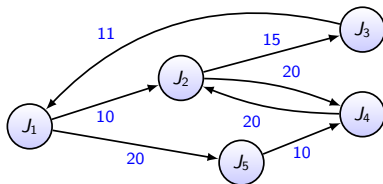


# Hierarchy of Models



# The Digraph Real-Time (DRT) Task Model

- Branching, cycles (loops), ...
- *Directed graph* for each task
  - ▶ Vertices  $J$ : jobs to be released (with WCET and deadline)
  - ▶ Edges  $(J_i, J_j)$ : minimum inter-release delays  $p(J_i, J_j)$

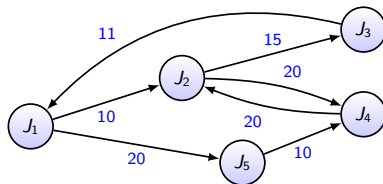


Theorem (S. et al., RTAS 2011)

For DRT task systems  $\tau$  with a utilization bounded by any  $c < 1$ , feasibility can be decided in *pseudo-polynomial time*.

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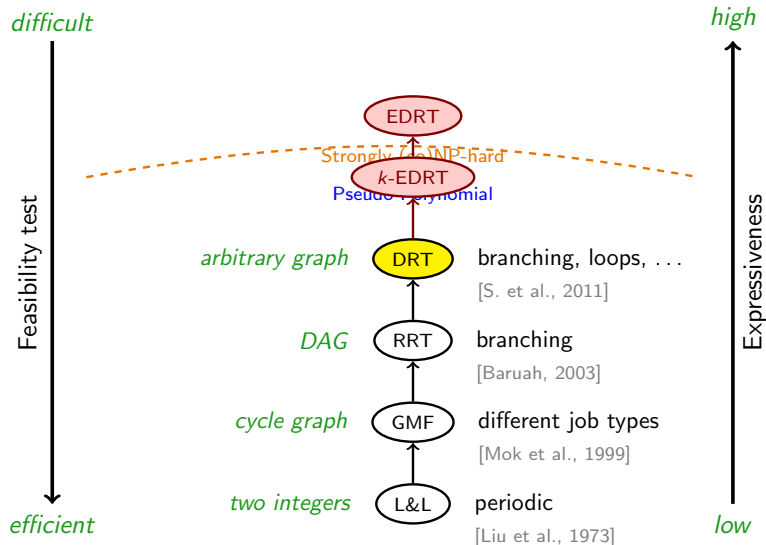
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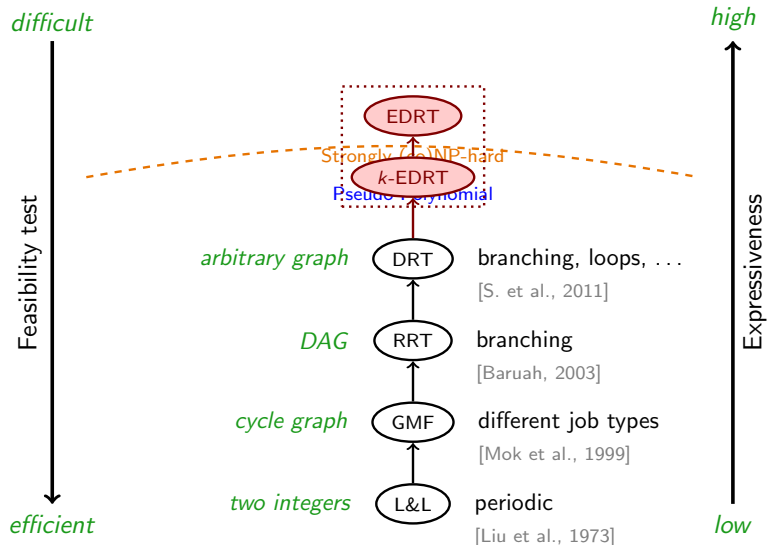
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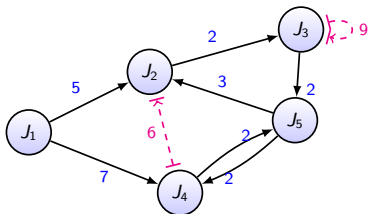
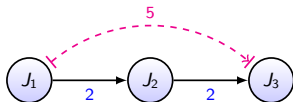


# Hierarchy of Models



# Extending DRT: Global Timing Constraints – *This Work*

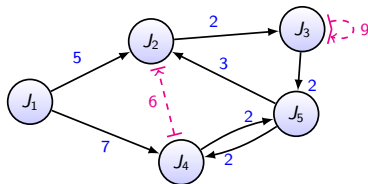
- Delays in DRT only for *adjacent* jobs
- What about adding *global delay constraints*?



- Motivation:
  - ▶ Mode sub-structures
  - ▶ Burstiness
  - ▶ ...

# Extended DRT (EDRT) – *This Work*

- Extends DRT with *global delay constraints*
- *Directed graph* for each task
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  - ▶  $k$  global constraints  $(J_i, J_j, \gamma)$



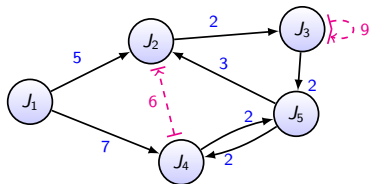
## Theorem (Our technical result)

For  $k$ -EDRT task systems with bounded utilization, feasibility is

- 1 decidable in *pseudo-polynomial time* if  $k$  is constant, and
- 2 *strongly coNP-hard* in general.

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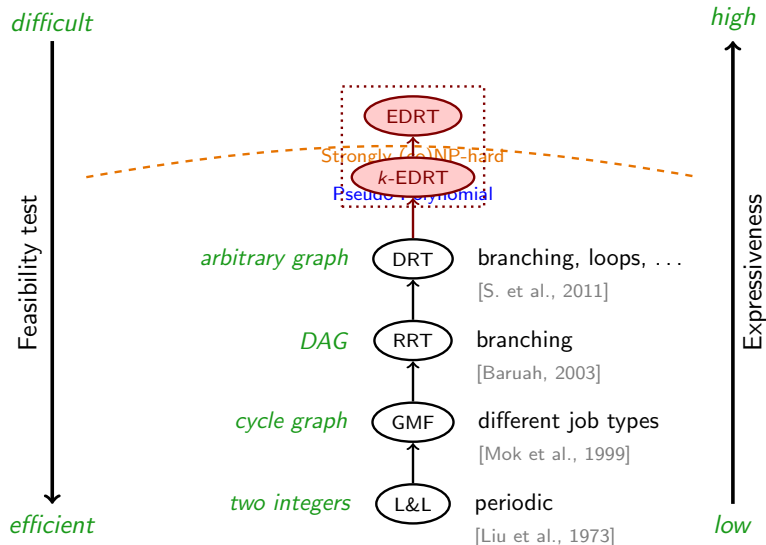
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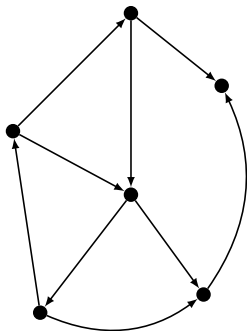


# Fahrplan

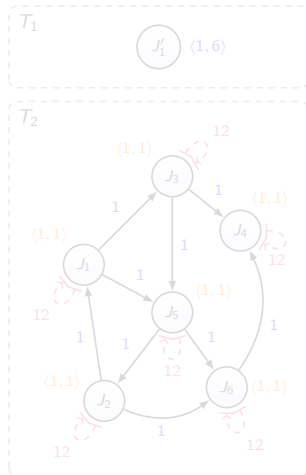
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# Hardness for EDRT

- Number of constraints now *not constant*
- Reduction from *Hamiltonian Path Problem* (strongly NP-hard)

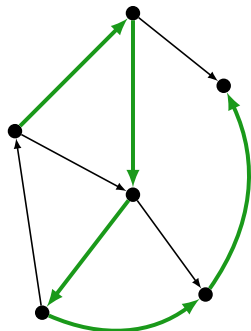


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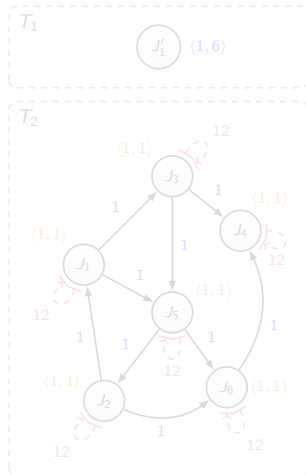


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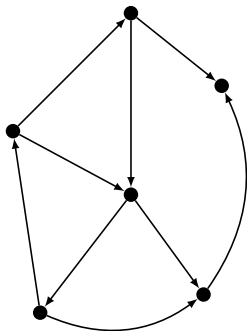


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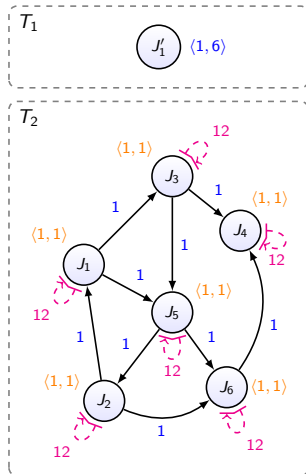


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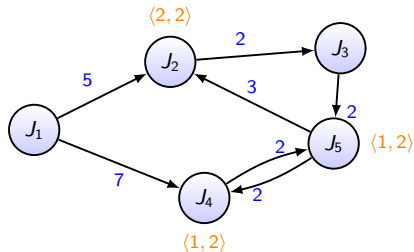


$\rightsquigarrow$



# Analysing $k$ -EDRT: The Problem

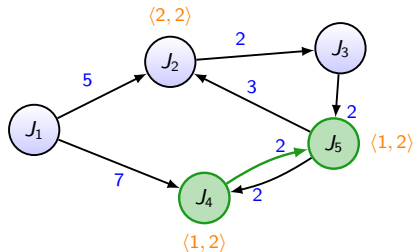
- Detour: Analyzing DRT
  - ▶ Using demand bound functions
  - ▶ Compute exec. demand and deadline for *all paths* in  $G(T)$



- Problem: Constraints ignored during path exploration
  - ▶  $\langle 2, 4 \rangle$  is lacking constraint information
  - ▶ ... about the *active constraint*

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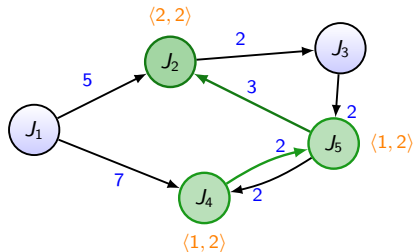
*Demand pair:  $\langle 2, 4 \rangle$*

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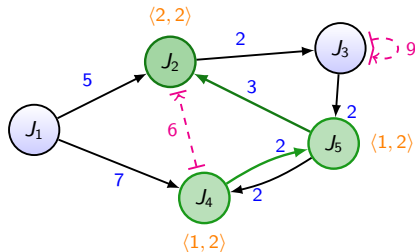
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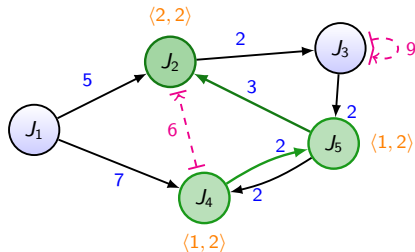
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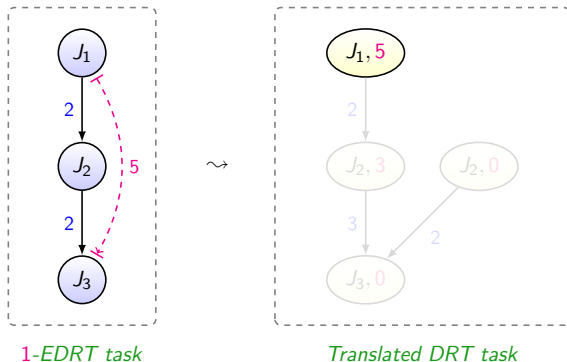
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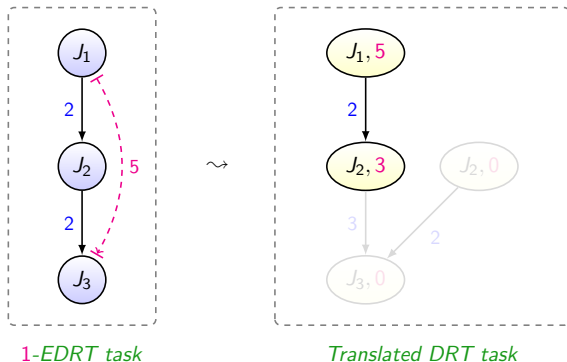
- Translate  $k$ -EDRT into *plain DRT*
  - ▶ Represent active constraints as *countdowns*
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- Optimizations: Efficient on-the-fly translation

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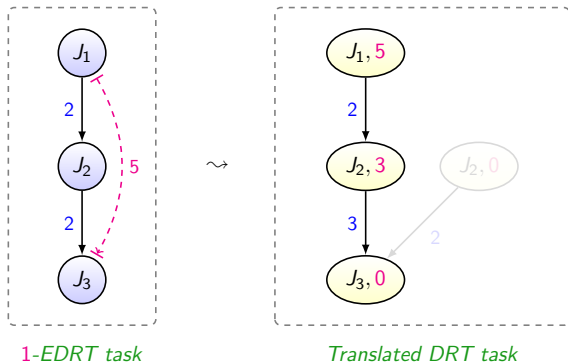
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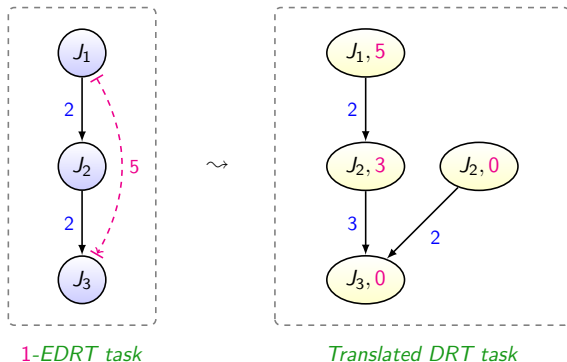
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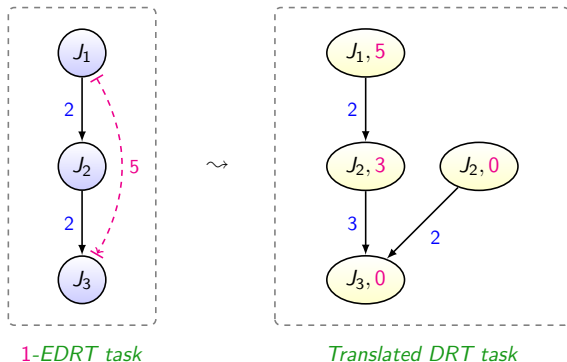
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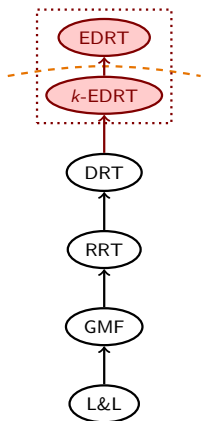


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# Summary and Outlook

- Introduced Extension to DRT Task Model
  - ▶ *Global delay constraints*
- Establishes tractability borderline for feasibility test
  - ▶ Constant number of constraints: *tractable*
  - ▶ Unbounded number of constraints: *intractable*
- Ongoing work:
  - ▶ Global constraints for simpler models (RRT, GMF)
  - ▶ Interaction with Resource Sharing Protocols (cf. talk tomorrow)



Thanks!