DATABASE TECHNOLOGY - 1DL116

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An introductory course on database systems

http://user.it.uu.se/~udbl/dbt-vt2007/
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Introduction to the Relational Model

Elmasri/Navathe ch 5, 7
Padron-McCarthy/Risch ch 5, 6

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The Relational Model

- The relational model was introduced by Dr. Edgar (Ted) F. Codd (1924-2003) in 1970.
  - Dr. Codd, a mathematician from Oxford (UK), was at that time working as an IBM researcher in the IBM San Jose Research Laboratory (USA).
- Many DBMS’s are based on the relational data model.
- It support simple declarative, but yet powerful, languages for describing operations on data.
- Operations in the relational model applies to relations (tables) and produce new relations.
  - This means that an operation can be applied to the result of another operation and that several different operations can be combined.
  - Operations are described in an algebraic notation that is based on relational algebra.
Relations as mathematical objects

- In set theory, a relation is defined as a subset of the product set (cartesian product) of a number of domains (value sets).
- The product set of the domains \(D_1, D_2, \ldots, D_n\) is written as \(D_1 \times D_2 \times \ldots \times D_n\).
- \(D_1 \times D_2 \times \ldots \times D_n\) constitute the set of all ordered sets \(<v_1, v_2, \ldots, v_n>\) such that \(v_i\) belongs to \(D_i\) for all \(i\).
  - If \(n=2\), \(D_1=\{T, F\}\) and \(D_2=\{P, Q, R\}\) one gets the product sets:
    - \(D_1 \times D_2 = \{<T,P>, <T,Q>, <T,R>, <F,P>, <F,Q>, <F,R>\}\)
    - \(D_2 \times D_1 = \{<P,T>, <P,F>, <Q,T>, <Q,F>, <R,T>, <R,F>\}\)
  - For example, we have the relations:
    - \(R_1 \subseteq D_2 \times D_1\) \(\Rightarrow R_1 = \{<P,T>, <Q,T>, <R,T>\}\)
    - \(R_2 \subseteq D_2 \times D_1\) \(\Rightarrow R_2 = \{<P,T>, <P,F>\}\)
- Members of a relation is called tuples. If the relation is of degree \(n\), the tuples are called \(n\)-tuples.
Relation schema and instance

- $A_1, A_2, \ldots, A_n$ are attributes
- $R = (A_1, A_2, \ldots, A_n)$ is a relation schema
  - Customer-schema(customer-name, customer-street, customer-city)
- $r(R)$ is a relation on the relation schema $R$
  - customer (Customer-schema)
- The current values (relation instance) of a relation are specified by a table.
- An element $t$ of $r$ is a tuple - represented by a row in a table customer

<table>
<thead>
<tr>
<th>customer-name</th>
<th>customer-street</th>
<th>customer-city</th>
</tr>
</thead>
<tbody>
<tr>
<td>Jones</td>
<td>Main</td>
<td>Harrison</td>
</tr>
<tr>
<td>Smith</td>
<td>North</td>
<td>Rye</td>
</tr>
<tr>
<td>Curry</td>
<td>North</td>
<td>Rye</td>
</tr>
<tr>
<td>Lindsay</td>
<td>Park</td>
<td>Pittsfield</td>
</tr>
</tbody>
</table>
First Normal Form

- Only simple or atomic values are allowed in the relational model.
- Attributes is not allowed to have composite or multiple values.
- The theory for the relational model is based on these assumptions which is called:

  *The first normal form assumption*
Null values

• A special value, null or ⊥, can sometimes be used as an attribute value.
• Every occurrence of null is unique. Thus, two occurrences of null is not considered to be equal even if they are represented by the same symbol.
• null is used:
  – when one does not know the actual value of an attribute.
  – when a certain attribute does not have a value.
  – when an attribute is not applicable.
• Examples of the use of null are showed later.
Keys

- Because relations are sets, all tuples in the relation are different.
- There is usually a subset $k$ of the attributes in a relation schema $R$, i.e. $k \subseteq R$, that has the characteristic that if the tuples $t_1, t_2 \in r(R)$ and $t_1 \neq t_2$, the following holds: $t_1[k] \neq t_2[k]$ (i.e. the value of $k$ in $t_1 \neq$ the value of $k$ in $t_2$)
- Every such subset $k$ is called a superkey for $R$. 
Keys - continued . . .

- A superkey $k$ is *minimal* if there is no other superkey $k'$ such that $k' \subset k$.

- Every minimal superkey (NOTE! there can be more than one) is called a **candidate key** for $R$.

- The candidate key chosen by the database designer as the key for $R$ is called $R$'s **primary key** or just **key**.

- In addition, term **foreign key** is used when a tuple is referenced, from another relation, with its key.
Determining keys from E-R types

• **Strong entity type.** The primary key of the entity type becomes the primary key of the relation.

• **Weak entity type.** The primary key of the relation consists of the union of the primary key of the strong entity type and the discriminator of the weak entity type.

• **Relationship type.** The union of the primary keys of the related entity types becomes a super key of the relation.
  – For binary many-to-many relationship types, above super key is also the primary key.
  – For binary many-to-one relationship types, the primary key of the “many” entity type becomes the relation’s primary key.
  – For one-to-one relationship types, the relation’s primary key can be that of either entity type.
Integrity constraints
for a relational database schema

• 1. Domain constraint
  – attribute values for attribute A shall be atomic values from \text{dom}(A)

• 2. Key constraint
  – candidate keys for a relation must be unique

• 3. Entity integrity constraint
  – no primary key is allowed to have a null value

• 4. Referential integrity constraint
  – a tuple that refers to another tuple in another relation must refer to an existing tuple

• 5. Semantic integrity constraint
  – e.g. “an employee’s total work time per week can not exceed 40 hours for all projects taken all together”
Steps in translation from E-R model to relational model

• Translation of entity types and their attributes
  – Step 1) Entity types
  – Step 2) Weak entity types

• Translation of relationships
  – Step 3) 1-1 Relationship
  – Step 4) 1-N Relationship
  – Step 5) M-N Relationship

• Translation of multivalued attributes and relationships
  – Step 6) Multivalued attributes
  – Step 7) Multivalued relationships
Translating entity types and their attributes

• Step 1: Entity types - a strong entity type reduces to a table with the same attributes.
  – Key attributes (primary key - pk) is made the primary key column(s) for the table. Each attribute gets their own column.
  – Composite attributes are normally represented by their simple components.
  – Example customer schema and table:

\[
\text{Customer(}\text{social-security, customer-name, c-street, c-city)}\]

<table>
<thead>
<tr>
<th>social-security</th>
<th>customer-name</th>
<th>c-street</th>
<th>c-city</th>
</tr>
</thead>
<tbody>
<tr>
<td>321-12-3123</td>
<td>Jones</td>
<td>Main</td>
<td>Harrison</td>
</tr>
<tr>
<td>019-28-3746</td>
<td>Smith</td>
<td>North</td>
<td>Rye</td>
</tr>
<tr>
<td>677-89-9011</td>
<td>Hayes</td>
<td>Main</td>
<td>Harrison</td>
</tr>
</tbody>
</table>
Translating entity types cont. . .

- **Step 2: Weak entity types** - a weak entity type becomes a table that includes a column for the primary key of the identifying strong entity type.

![Diagram of entity types](image)
Translating entity types cont. . .

- The table corresponding to a relationship type linking a weak entity type to its identifying strong entity type is redundant.

- Example of the payment schema and table:
  - The payment table already contains the information that would appear in the loan-payment table (i.e., the columns loan-number and payment-no).

\textit{Payment}(\texttt{loan-number, payment-no, pay-date, amount})

<table>
<thead>
<tr>
<th>loan-number</th>
<th>payment-no</th>
<th>pay-date</th>
<th>amount</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-17</td>
<td>5</td>
<td>10 May 1996</td>
<td>50</td>
</tr>
<tr>
<td>L-23</td>
<td>11</td>
<td>17 May 1996</td>
<td>75</td>
</tr>
<tr>
<td>L-15</td>
<td>22</td>
<td>23 May 1996</td>
<td>300</td>
</tr>
</tbody>
</table>
Translating relationship types

- Step 3: 1-1 Relationship types
  - The foreign key column (fk) is a copy of the other entity’s primary key column (pk). The values in a fk-column point to unique row in the other table, and thus implement the relationship.

```
Alt 1:

<table>
<thead>
<tr>
<th>pk1</th>
<th>a1</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>pk2</th>
<th>a2</th>
<th>fk1</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Alt 2:

<table>
<thead>
<tr>
<th>pk1</th>
<th>a1</th>
<th>fk2</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>pk2</th>
<th>a2</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
```
Translating 1-1 relationship types cont...
Translating relationship . . . cont.

• Step 4: 1-N Relationship types
  – Include the primary key of the “1-side” as a foreign key on the “N-side”, (i.e. the foreign key column is placed on the entity on the N-side).
  – Alternatively, an extra table (R) is created whose primary key is a foreign key composed by the primary key from the N-side.
Translating relationship . . . cont. . .

- Step 5: M-N Relationship types
  - Always a separate table with columns for the primary keys of the two participating entity types, and any descriptive attributes of the relationship type.
Translating relationship ... cont.

- **Step 6: Multivalued attributes**
  - A separate table is created for the multivalued attribute. Its primary key is composed of the owning entity’s primary key, and the attribute value itself.
Translating relationship . . . cont.

- **Step 7: Multivalued relationship types**
  - First try to remove multivalued relationships on the E-R model level by model transformation.
  - A separate table is created, with foreign keys to all tables that are included in the relationship. Its primary key is composed of all foreign keys.

![Diagram](image-url)
Translating relationship . . . cont.

- Step 7: Multivalued relationship types continued
  - In the case where R is 1-N-N, the primary key on R shall not include the fk for the table with cardinality 1.
Translating Specialization/Generalization

• Alternative a) in Elmasri/Navathe

```
E1
  pk
a1
E2
  pk
a2
E3
  pk
a3
```

E1

pk

E2

pk

E3

pk

Translating aggregation

- Translating an implicit aggregation relationship type.

![Diagram showing part_of relationship]

- Translating an objectified aggregation relationship type.

![Diagram showing consist and part_in relationships]
Example E-R to relational model translation

- **EMPLOYEE**
  - `ename`
  - `salary`
  - **WORKS_IN**
    - `dno`
    - `dname`
  - **MANAGES**
    - `dno`
    - `dname`

- **DEPARTMENT**
  - `dno`
  - `dname`
  - **CARRIES**
    - `dno`
    - `sno`
    - `sname`
    - `saddr`

- **ITEM**
  - `iname`
  - `ino`
  - **INCLUDE**
    - `oni`
    - `quantity`

- **ORDER**
  - `ono`
  - **PLACED_BY**
    - `ono`
    - `cname`
    - `caddr`
    - `balance`

- **CUSTOMER**
  - `cname`
  - `caddr`
  - `balance`

- **SUPPLIER**
  - `sname`
  - `saddr`
  - **SUPPLIES**
    - `sno`
    - `price`

- **INCLUDE**
  - `price`
  - `quantity`