

# Process Management

## Process Concepts

Frédéric Haziza <daz@it.uu.se>

Department of Computer Systems  
Uppsala University

Spring 2007

# Outline

- 1 Introduction
  - Definition
  - States
  - PCB
  - Transitions
- 2 Scheduling
- 3 Communication
  - IPC
  - Models
- 4 Threads
  - From process flows
  - Benefits
  - Models
  - Issues

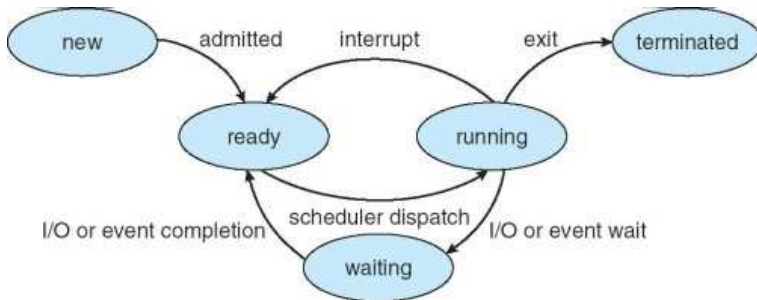


# What characterizes a process?

- Program in execution
- Stack (Temporary data, function parameters,...)
- Heap
- Data section (Global variables)
- CPU Registers
- Program Counter (PC)
  
- Program code = Text section
- Program in execution = text section (executable file) *loaded in memory*



# States



# Process Control Block (PCB)

<b>PCB</b>
process state
process ID (number)
PC
Registers
memory information
open files
other resources

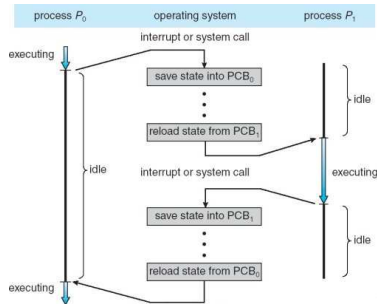


- Interrupts

- Generated asynchronously by external devices and timers
- Example: The I/O is complete or timers have expired

- Traps (software errors, illegal instructions)

- System calls



# Scheduling Queues

## Job Queue

Linked list of PCBs

- (main) job queue
- ready queue
- device queues



# Scheduling

- **Job scheduler** (loads from disk)
- **CPU scheduler** (dispatches from ready queue)





# Context Switch

## Context switch

### PCB swap

- Cost?
- 10ms switch for 100ms work => 9% wasted



# Process creation and termination

See the lab 1 ...



# Interprocess Communication (IPC)

## 2 models

- Message Passing
- Shared Memory

See black board...

## Benefits

- Small amount to exchange  
=> Message Passing, because no conflict to avoid
- Shared Memory  
=> Working at the speed of memory – faster



## Shared Memory

Recall that the OS prevents processes to share memory => Agreement on relaxing restriction

### Example (Producer-Consumer)

Unbounded buffer and bounded buffer  
(book p98)

Requires:

- Synchronisation  
(No consumption of non produced items)
- Waiting

## Message Passing

No shared space.  
Can be distributed across network

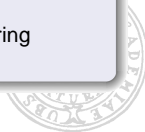
### Example

Chat program

- send(m)
- receive(m)

Requires a communication link

- direct or indirect  
(mailbox/ports)
- synch. or asynch.  
(blocking or non-blocking)
- automatic or explicit buffering  
(info on the link)



# Threads

## Heavy-weight vs Light-weight...

### Example (Web server)

We want to serve more than one client at a time

- 1 process. If incoming request, new process created => costly!
- 1 process. If same task as other one, why overhead => better to multithread

On Solaris:

- Time for creating a process = 30 x time for creating a thread
- Time for context switching = 5 x time for switching a thread



# Benefits

- Responsiveness
- Resource sharing
- Economy
- Utilization of multiprocessor architectures



# Multithread Models

Deals with correspondance between

- threads in **user space**
- threads in **kernel space**

**One to One**

**Many to One**

**Many to Many**



# Issues

- *fork()* creates a copy of a process with all threads  
or  
just the one which calls the fork?
- Cancellation
- Signal handling: Read in book page 139
- Thread pool (limit in system, pre-create threads)
- Thread specific data (sharing data?)
- about the models themselves

