

# UPPSALA UNIVERSITET

TOWARD A UNIFIED TASK-BASED PARALLEL PROGRAMMING INTERFACE (UTP) Afshin Zafari, Elisabeth Larsson IT Department, Division of Scientific Computing

#### A sample GEMM program

#### int main(int argc, char \*\*argv)

utp\_initialize(argc,argv);

// UTP start

int M = config.getYDimension(); // Get parameters int B1 = config.getYBlocks(1); int B2 = config.getYBlocks(2); GData A(M,M),B(M,M),C(M,M); // Define Data GPartitioner P1(B1,B1); GPartitioner P2(B2,B2); // Define Partitions P1->set\_next(P2); // Define Partitions A.set\_partition(P1); // Apply partitioninng on data B.set\_partition(P1);

## What is it?

- Provides a **unified** programming interface for task based parallel programming.
- Enables **different task-based frameworks** to cooperate without knowing each other.
- Enables a program written **once** in **sequential** form to run in **parallel** on different types of computing resources.

### Task and Data Hierarchy



<pre>C.set_partition(P1);</pre>	
ugemm(A,B,C);	

utp\_finalize();

// UTP waits for all tasks

#### ugemm

```
void ugemm(GData &A,GData &B,GData &C,GTask *p=NULL){
    int m = A.get_part_countY();
    int n = B.get_part_countX();
    int o = C.get_part_countX();
    for (int i = 0; i < m; i++)
        for (int j = 0; j < n; j++)
            for (int k = 0; k < o; k++)
               ugemm_t(A(i,k),B(k,j),C(i,j),p);
}
void ugemm_t(GData &A,GData &B,GData &C,GTask *p){
        packArgs( args, A , B , C );
        packAxs ( axs , In , In , InOut );
        get_dispatcher()->submit_task(ugemmo,args,axs,p);
}
void ugemmo::split(GTask *t){
```

```
//unpack arguments of t to A,B,C: A = t->args[0] ,...
ugemm(A,B,C,t);
```

## Which scheduler or framework?

### Generic Objects

#### • Generic Data

- M, N, memory, lead\_dim, partition, parent, level
- Generic Partition
- MB, NB, parent, get\_part(i,j)
- Generic Task
   Operation, parent, args[], access[],
   level, kernel
- Generic Operation  $f(x, y) \mapsto \{ tasks(x_{ij}, y_{k\ell}) \}$
- Generic Scheduler submit(), run(), finished() Notifications (ready and finished)
- **Dispatch Policy** Used for customizing the chain of schedulers
- Dispatcher
  - submit(), run(), finished()

#### $t_m^{\ell}$ : task m at level $\ell$ $x_{i,jk}$ : partition jk of $x_i$

### Where do the tasks run?

The configuration of computing resources can be determined at run time.

#### run ugemm\_app

--cores <multi-core configuration>
--gpus <gpu configuration>
--nodes <dist. mem. configuration>
--cloud <cloud configuration>
--cluster <cluster configuration>

#### Configuration Sample

	Application	programs
$\frown$	💻 that use U	nified Task



- Wrappers translate framework interfaces to the unified interface.
  BLAS and cuBLAS wrappers used as schedulers that run kernels.
  Message Queue (MQ) Scheduler
  puts the commands received from Dispatcher to a public queue.
- -gets messages from a public queue and sends them to Dispatcher.

## More information?

Central hub of *schedulers* conversations.

does it work? How **Schedulers** Dispatch Data, Dispatcher Program Operation Policy Partitions  $f(x_1, x_2, ...)$ submit submit tasks at  $\ell$ =1 ready to run split submit tasks at  $\ell$ =2 submit finished finished

- -Resulting tasks of calling an Operation are submitted to Schedulers via Dispatcher.
- -Ready task at any level of the hierarchy splits
- smog@cloud.snic.se Interface can run on cluster, cloud and client Using MQ computers. DuctTei ouctTeip Scheduler, all instances StarPU of UTP programs can BLAS communicate tasks and data with each other. Why to use it? • Independent from frameworks
- *Transparent* to new technologies in the underlying hardware
- Decoupled application programming
- Mixing different frameworks
- Single application program for any available parallelism
- New features in any framework become available





